Sorting Algorithms

Lecture 03

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- Inroduction
 - Comparision
- Bubble Sort
 - Algo
 - Example
 - Analysis
- Insertion Sort
 - Algo
 - Example
 - Analysis
- Selection Sort
 - Algo
 - Example
 - Analysis
- Merge Sort
 - Algo
 - Example
 - Analysis

Outline

- Inroduction
 - Comparision
- **Bubble Sort**
 - Algo

 - Analysis
- Insertion Sort
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 - Analysis
- - Algo
 - Example
 - Analysis
- Merge Sort
 - Algo



Problem

input: A sequence of *n* numbers $< a_1, a_2, \ldots, a_n >$. **output:** A permutation (reordering) $< a'_1, a'_2, \ldots, a'_n >$ of the input sequence such that $a'_1 \le a'_2 \le \ldots \le a'_n$.

Example

$$<$$
 31, 41, 59, 26, 41, 58 $>\Rightarrow$ $\boxed{\mathsf{SORTING}}$ \Rightarrow $<$ 26, 31, 41, 41, 58, 59 $>$

Applications

- Easy search
- Database operations
- Basis of many algos (string processing, partitioning, intersection)

Factors

- Running-time (worst, best, average): Asymptotic time-complexity
- Additional space requirement: In-place ($\theta(1)$ memory) or not
- Initial conditions(input order, key distribution, size) of input data: already-sorted, nearly-sorted, reverse-sorted, few unique keys, small/large size
- Stability: => two elements which have the same value will retain their relative order after sorting. Example:

$$<$$
 31, 41*, 41+, 26, 41-, 58 $>\Rightarrow$ SORTING \Rightarrow $<$ 26, 31, 41*, 41+, 41-, 58 $>$



An ideal algorithm

- O(nlogn) worst time comparisons $(\Omega(n))$ lower bound with certain assumption about the data. In general can not be less than this a.
- In-place
- Stable
- Adaptive

No algorithm is an absolute ideal.



^aA more detailed proof is given in Donald E. Knuth, The Art of Computer Programming, Volume 3: Sorting and Searching, 2nd Ed., Addison Wesley, 1998, §5.3.1, p.180.)

Types

- External
- Internal
 - Bubble
 - Insertion
 - Selection
 - Merge

Outline

- - Comparision
- **Bubble Sort**
 - Algo

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Pass 1

_ _ _ _ _ _



. . .



5 1 2 0 4 6 0 9

5 1 2 0 4 6 0 9

Pass 1

5 1 2 0 4 6 0 9

5 1 2 0 4 6 0 9

Pass 1

1 5 2 0 4 6 0 9

5 1 2 0 4 6 0 9

Pass 1

1 2 5 0 4 6 0 9

5 1 2 0 4 6 0 9

Pass 1

1 2 0 5 4 6 0 9

5 1 2 0 4 6 0 9

Pass 1

1 2 0 4 5 6 0 9

5 1 2 0 4 6 0 9

Pass 1

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Pass 1

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Pass 1



5 1 2 0 4 6 0 9

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Pass 1

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Pass 2

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5 1 2 0 4 6 0 9

Pass 1

1 2 0 4 5 0 6 9

Pass 2

10240569

5 1 2 0 4 6 0 9

Pass 1

1 2 0 4 5 0 6 9

Pass 2

10240569

5 1 2 0 4 6 0 9

Pass 1

1 2 0 4 5 0 6 9

Pass 2

1 0 2 4 0 5 6 9

Pass 3

1 0 2 4 0 5 <u>6</u> 9

5 1 2 0 4 6 0 9

Pass 1

1 2 0 4 5 0 6 9

Pass 2

1 0 2 4 0 5 6 9

Pass 3

0 1 2 4 0 5 6 9

5 1 2 0 4 6 0 9

Pass 1

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Pass 4

Input

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Pass 7

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Pass 2

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Pass 4

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Pass 6

5 1 2 0 4 6 0 9

Pass 1

1 2 0 4 5 0 6 9

Pass 3

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Pass 5

0 0 1 2 4 5 6

Pass 7

0 0 1 2 4 5 6 9

Pass 2

10240569

Pass 4

0 1 0 2 4 5 6 9

Pass 6

5 1 2 0 4 6 0 9

Output





Algorithm 1 BubbleSort algorithm

```
1: procedure BUBBLE(a[], n)
       for i \leftarrow n-1 to 1 do
2:
3:
           for i \leftarrow 1 to i do
              if a[i - 1] > a[j] then
4:
                  swap(a[i - 1], a[i])
5:
              end if
6:
           end for
7:
8:
       end for
9: end procedure
```

Algorithm 2 BubbleSort algorithm

```
1: procedure BUBBLE(a[], n)
      for i \leftarrow n-1 to 1 do
2:
          for i \leftarrow 1 to i do
3:
              if a[i - 1] > a[i] then
4:
                  swap(a[j-1],a[j])
5:
              end if
6:
          end for
7:
      end for
8:
9: end procedure
```

Algorithm 3 BubbleSort algorithm

```
1: procedure BUBBLE(a[], n)
      for i \leftarrow n-1 to 1 do
2:
          for i \leftarrow 1 to i do
3:
              if a[i - 1] > a[i] then
4:
                  swap(a[j-1],a[j])
5:
              end if
6:
          end for
7:
      end for
8:
9: end procedure
```

Algorithm 4 BubbleSort algorithm

```
1: procedure BUBBLE(a[], n)
      for i \leftarrow n-1 to 1 do
2:
          for i \leftarrow 1 to i do
3:
              if a[i - 1] > a[i] then
4:
                  swap(a[j-1], a[j])
5:
              end if
6:
          end for
7:
      end for
8:
9: end procedure
```

Algorithm 5 BubbleSort algorithm

```
1: procedure BUBBLE(a[], n)
      for i \leftarrow n-1 to 1 do
          for i \leftarrow 1 to i do
3:
              if a[i - 1] > a[i] then
4:
                  swap(a[j-1],a[j])
5:
              end if
6:
          end for
7:
      end for
8:
9: end procedure
```

Running Time

$$\Sigma_0^{n-1}\Theta(i) = \Theta(\Sigma_1^{n-1}i) = \Theta(n^2)$$

- Worst case:
- Reverse sorted:
- Already sorted:

In-place

Stable

Adaptive

- Worst case: $O(n^2)$
- Reverse sorted:
- Already sorted:

In-place

Stable

Adaptive

- Worst case: O(n²)
- Reverse sorted: $\Theta(n^2)$
- Already sorted:

In-place

Stable

Adaptive

- Worst case: O(n²)
- Reverse sorted: $\Theta(n^2)$
- Already sorted: Can be designed for $\Theta(n)$

In-place

Stable

Adaptive



- Worst case: O(n²)
- Reverse sorted: $\Theta(n^2)$
- Already sorted: Can be designed for $\Theta(n)$

In-place

Stable

Adaptive

- Worst case: $O(n^2)$
- Reverse sorted: $\Theta(n^2)$
- Already sorted: Can be designed for $\Theta(n)$

In-place

O(1) extra space

Stable

Adaptive



- Worst case: O(n²)
- Reverse sorted: $\Theta(n^2)$
- Already sorted: Can be designed for $\Theta(n)$

In-place

O(1) extra space

Stable

Yes

Adaptive



- Worst case: O(n²)
- Reverse sorted: $\Theta(n^2)$
- Already sorted: Can be designed for $\Theta(n)$

In-place

O(1) extra space

Stable

Yes

Adaptive



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An intermediate step

• • • • • • •

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• • • • • • •

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Pass 1

_ _ _ _ _ _ _ _ _ _

• • • • • • •

• • • • • • • • •

Pass 2

. . .

Pass 7

. **.**

5 1 2 0 4 6 0 9

5 1 2 0 4 6 0 9

Pass 1

5 1 2 0 4 6 0 9

5 1 2 0 4 6 0 9

Pass 1

1 5 2 0 4 6 0 9

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Pass 1

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Pass 1

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Pass 7

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Pass 3

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Pass 5

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Pass 2

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Pass 4

0 1 2 4 5 6 0 9

Pass 6

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Pass 7

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5 1 2 0 4 6 0 9

Pass 1

1 5 2 0 4 6 0 9

Pass 3

0 1 2 5 4 6 0 9

Pass 5

0 1 2 4 5 6 0 9

Pass 2

1 2 5 0 4 6 0 9

Pass 4

0 1 2 4 5 6 0 9

Pass 6

0 0 1 2 4 5 6 9

Pass 7

0 0 1 2 4 5 6 9

5 1 2 0 4 6 0 9

Output





Algorithm 6 Insertion algorithm

```
1: procedure INERTION(a[], n)
        for i \leftarrow 1 to n-1 do
 2:
             key \leftarrow a[i]
 3:
            i \leftarrow i
 4:
             while j > 0 and a[j-1] > key do
 5:
                 a[i] \leftarrow a[i-1]
 6:
                 j \leftarrow j - 1
 7:
             end while
 8:
             a[j] \leftarrow key
 9:
        end for
10:
11: end procedure
```

Algorithm 7 Insertion algorithm

```
1: procedure INERTION(a[], n)
        for i \leftarrow 1 to n-1 do
             key \leftarrow a[i]
 3:
            i \leftarrow i
 4:
            while j > 0 and a[j-1] > key do
 5:
                 a[j] \leftarrow a[j-1]
 6:
                 j \leftarrow j - 1
 7:
             end while
 8:
 9:
             a[i] \leftarrow key
        end for
10:
11: end procedure
```

Algorithm 8 Insertion algorithm

```
1: procedure INERTION(a[], n)
        for i \leftarrow 1 to n-1 do
             key \leftarrow a[i]
 3:
            i \leftarrow i
 4:
            while j > 0 and a[j-1] > key do
 5:
                 a[j] \leftarrow a[j-1]
 6:
                 j \leftarrow j - 1
 7:
             end while
 8:
 9:
             a[i] \leftarrow key
        end for
10:
11: end procedure
```

Algorithm 9 Insertion algorithm

```
1: procedure INERTION(a[], n)
         for i \leftarrow 1 to n-1 do
             key \leftarrow a[i]
 3:
             i \leftarrow i
 4:
             while j > 0 and a[j-1] > key do
 5:
                 a[j] \leftarrow a[j-1]
 6:
                 j \leftarrow j - 1
 7:
             end while
 8:
 9:
             a[i] \leftarrow key
         end for
10:
```

11: end procedure

Algorithm 10 Insertion algorithm

```
1: procedure INERTION(a[], n)
        for i \leftarrow 1 to n-1 do
             key \leftarrow a[i]
 3:
            i \leftarrow i
            while j > 0 and a[j-1] > key do
 5:
                 a[j] \leftarrow a[j-1]
                 j \leftarrow j - 1
 7:
             end while
 8:
 9:
             a[i] \leftarrow key
        end for
10:
11: end procedure
```

Running Time

 $\sum_{i=1}^{n-1} t_i$ where t_i is the time taken in i^{th} loop.

- Worst case:
- Best case:
- Reverse sorted:
- Already sorted:

In-place

Stable

- Worst case: $t_i = i \Rightarrow \text{Running time } \Sigma_1^{n-1} i = O(n^2)$
- Best case: $t_i = 1 \Rightarrow \text{Running time } \Sigma_1^{n-1} 1 = O(n)$
- Reverse sorted:
- Already sorted:

In-place

Stable



- Worst case: $t_i = i \Rightarrow \text{Running time } \Sigma_1^{n-1} i = O(n^2)$
- Best case: $t_i = 1 \Rightarrow \text{Running time } \Sigma_1^{n-1} 1 = O(n)$
- Reverse sorted: $\Theta(n^2)$
- Already sorted:

In-place

Stable



- Worst case: $t_i = i \Rightarrow \text{Running time } \Sigma_1^{n-1} i = O(n^2)$
- Best case: $t_i = 1 \Rightarrow \text{Running time } \Sigma_1^{n-1} 1 = O(n)$
- Reverse sorted: $\Theta(n^2)$
- Already sorted: $\Theta(n)$

In-place

Stable



- Worst case: $t_i = i \Rightarrow \text{Running time } \Sigma_1^{n-1} i = O(n^2)$
- Best case: $t_i = 1 \Rightarrow \text{Running time } \Sigma_1^{n-1} 1 = O(n)$
- Reverse sorted: $\Theta(n^2)$
- Already sorted: $\Theta(n)$

In-place

Stable



- Worst case: $t_i = i \Rightarrow \text{Running time } \Sigma_1^{n-1} i = O(n^2)$
- Best case: $t_i = 1 \Rightarrow \text{Running time } \Sigma_1^{n-1} 1 = O(n)$
- Reverse sorted: $\Theta(n^2)$
- Already sorted: $\Theta(n)$

In-place

O(1) extra space

Stable



- Worst case: $t_i = i \Rightarrow$ Running time $\Sigma_1^{n-1}i = O(n^2)$
- Best case: $t_i = 1 \Rightarrow \text{Running time } \Sigma_1^{n-1} 1 = O(n)$
- Reverse sorted: $\Theta(n^2)$
- Already sorted: $\Theta(n)$

In-place

O(1) extra space

Stable

Yes



- Worst case: $t_i = i \Rightarrow$ Running time $\Sigma_1^{n-1}i = O(n^2)$
- Best case: $t_i = 1 \Rightarrow \text{Running time } \Sigma_1^{n-1} 1 = O(n)$
- Reverse sorted: $\Theta(n^2)$
- Already sorted: $\Theta(n)$

In-place

O(1) extra space

Stable

Yes

Adaptive

Yes, and less overhead



- Worst case:
- Best case:
- Reverse sorted:
- Already sorted:

In-place

Stable

Adaptive

What if binary search is used?

Running Time

- Worst case:
- Best case:
- Reverse sorted:
- Already sorted:

In-place

Stable

Adaptive

What if a linked list is used instead of an array?

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An intermediate step





• •

Pass 2





. .

Pass 1





Pass 7







5 1 2 0 4 6 0 9

Pass 1

5 1 2 0 4 6 0 9

Pass 1

5 1 2 0 4 6 0 9

Pass 1

5 1 2 0 4 6 0 9

Pass 1

5 1 2 0 4 6 0 9

Pass 1

0 1 2 5 4 6 0 9

Pass 2

5 1 2 0 4 6 0 9

Pass 1

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Pass 1

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Pass 3

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Pass 4

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Pass 1

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Pass 3

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Pass 2

0 0 2 5 4 6 1 9

Pass 4

0 1 5 4 6 2 9

5 1 2 0 4 6 0 9

Pass 1

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Pass 3

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Pass 2

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Pass 4

5 1 2 0 4 6 0 9

Pass 1

0 1 2 5 4 6 0 9

Pass 3

0 0 1 5 4 6 2 9

Pass 2

0 0 2 5 4 6 1 9

Pass 4

0 0 1 2 4 6 5 9

Pass 5

5 1 2 0 4 6 0 9

Pass 1

0 1 2 5 4 6 0 9

Pass 3

0 0 1 5 4 6 2 9

Pass 2

0 0 2 5 4 6 1 9

Pass 4

0 0 1 2 4 6 5 9

Pass 5

5 1 2 0 4 6 0 9

Pass 1

0 1 2 5 4 6 0 9

Pass 3

0 0 1 5 4 6 2 9

Pass 2

0 0 2 5 4 6 1 9

Pass 4

0 0 1 2 4 6 5 9

Pass 5

5 1 2 0 4 6 0 9

Pass 1

0 1 2 5 4 6 0 9

Pass 3

0 0 1 5 4 6 2 9

Pass 2

0 0 2 5 4 6 1 9

Pass 4

0 0 1 2 4 6 5 9

Pass 5

5 1 2 0 4 6 0 9

Pass 1

0 1 2 5 4 6 0 9

Pass 3

0 0 1 5 4 6 2 9

Pass 5

0 0 1 2 4 6 5 9

Pass 2

0 0 2 5 4 6 1 9

Pass 4

0 0 1 2 4 6 5 9

Pass 6

5 1 2 0 4 6 0 9

Pass 1

0 1 2 5 4 6 0 9

Pass 3

0 0 1 5 4 6 2 9

Pass 5

0 0 1 2 4 6 5 9

Pass 2

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Pass 4

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Pass 6

5 1 2 0 4 6 0 9

Pass 1

0 1 2 5 4 6 0 9

Pass 3

0 0 1 5 4 6 2 9

Pass 5

0 0 1 2 4 6 5 9

Pass 2

0 0 2 5 4 6 1 9

Pass 4

0 0 1 2 4 6 5 9

Pass 6

5 1 2 0 4 6 0 9

Pass 1

0 1 2 5 4 6 0 9

Pass 3

0 0 1 5 4 6 2 9

Pass 5

0 0 1 2 4 6 5 9

Pass 2

0 0 2 5 4 6 1 9

Pass 4

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Pass 6

0 0 1 2 4 5 6 9

5 1 2 0 4 6 0 9

Pass 1

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Pass 3

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Pass 5

0 0 1 2 4 6 5 9

Pass 2

0 0 2 5 4 6 1 9

Pass 4

0 0 1 2 4 6 5 9

Pass 6

0 0 1 2 4 5 6 9

Pass 7

0 0 1 2 4 5 6 9

2 0 4 6 0 9

Pass 1

5 4 6 0 9

Pass 3

5 4 6 2 9

Pass 5

6 5 9

Pass 7

5 6

Pass 2

5 4 6

Pass 4

6 5 9 4

Pass 6

6 9 5

Pass 2

Pass 4

Input

5 1 2 0 4 6 0 9

Pass 1

0 1 2 5 4 6 0 9

Pass 3

0 0 1 5 4 6 2 9

Pass 5

0 0 1 2 4 6 5 9

ol a lo

0 0 1 2 4 6 5 9

5 4 6

Pass 6

0 0 1 2 4 5 6 9

Pass 7

0 0 1 2 4 5 6 9

5 1 2 0 4 6 0 9

Output





Algorithm 11 Selection algorithm

```
1: procedure SELECTION(a[], n)
        for i \leftarrow 0 to n-2 do
 2:
            minPos \leftarrow i
 3:
            for j \leftarrow i + 1 to n - 1 do
 4:
                if a[j] < a[minPos] then
 5:
                    minPos \leftarrow i
 6:
 7:
                end if
            end for
 8:
            swap(a[i], a[minPos])
 9:
        end for
10:
11: end procedure
```

Algorithm 12 Selection algorithm

```
1: procedure SELECTION(a[], n)
        for i \leftarrow 0 to n-2 do
            minPos \leftarrow i
 3:
            for j \leftarrow i + 1 to n - 1 do
 4:
               if a[j] < a[minPos] then
 5:
                    minPos ← j
 6:
                                               \Theta(1)
               end if
 7:
            end for
 8:
            swap(a[i], a[minPos])
 9:
        end for
10:
11: end procedure
```

Algorithm 13 Selection algorithm

```
1: procedure SELECTION(a[], n)
      for i \leftarrow 0 to n-2 do
         minPos \leftarrow i
3:
         for j \leftarrow i + 1 to n - 1 do
4:
           5:
6:
7:
         end for
8:
         swap(a[i], a[minPos])
9:
      end for
10:
11: end procedure
```

Algorithm 14 Selection algorithm

```
1: procedure SELECTION(a[], n)
        for i \leftarrow 0 to n-2 do
             minPos \leftarrow i
 3:
             for j \leftarrow i + 1 to n - 1 do
 4:
                 if a[j] < a[minPos] then
 5:
                                                  \Theta(1) \Theta(n-i) \Sigma_0^{n-2}\Theta(n-i) times
                     minPos \leftarrow i
 6:
                 end if
 7:
             end for
 8:
             swap(a[i], a[minPos])
 9:
        end for
10:
11: end procedure
```

Algorithm 15 Selection algorithm

```
1: procedure SELECTION(a[], n)
         for i \leftarrow 0 to n-2 do
              minPos \leftarrow i
 3:
              for j \leftarrow i + 1 to n - 1 do
 4:
                   if a[j] < a[minPos] then
 5:
                                                      \Theta(1) \begin{cases} \Theta(n-i) \\ \text{times} \end{cases} \sum_{0}^{n-2} \Theta(n-i)
                        minPos \leftarrow i
 6:
                   end if
 7:
              end for
 8:
              swap(a[i], a[minPos])
 9:
         end for
10:
11: end procedure
```

Running Time

$$\Sigma_0^{n-2}\Theta(n-i) = \Theta(\Sigma_1^{n-1}(n-i)) = \cdots = \Theta(\Sigma_1^n i) = \Theta(n^2)$$

- Worst case:
- Best case:
- Reverse sorted:
- Already sorted:

In-place

Stable

- Worst case: $O(n^2)$
- Best case: $O(n^2)$
- Reverse sorted:
- Already sorted:

In-place

Stable

- Worst case: $O(n^2)$
- Best case: $O(n^2)$
- Reverse sorted: $\Theta(n^2)$
- Already sorted:

In-place

Stable

• Worst case: $O(n^2)$

• Best case: $O(n^2)$

• Reverse sorted: $\Theta(n^2)$

• Already sorted: $\Theta(n^2)$

In-place

Stable

• Worst case: $O(n^2)$

• Best case: $O(n^2)$

• Reverse sorted: $\Theta(n^2)$

• Already sorted: $\Theta(n^2)$

In-place

Stable

- Worst case: $O(n^2)$
- Best case: $O(n^2)$
- Reverse sorted: $\Theta(n^2)$
- Already sorted: $\Theta(n^2)$

In-place

O(1) extra space

Stable



- Worst case: $O(n^2)$
- Best case: $O(n^2)$
- Reverse sorted: $\Theta(n^2)$
- Already sorted: $\Theta(n^2)$

In-place

O(1) extra space

Stable

No



- Worst case: $O(n^2)$
- Best case: $O(n^2)$
- Reverse sorted: $\Theta(n^2)$
- Already sorted: $\Theta(n^2)$

In-place

O(1) extra space

Stable

No

Adaptive

NO

- Worst case:
- Best case:
- Reverse sorted:
- Already sorted:

In-place

Stable

Adaptive

Is there anything good about this algo? What if a heap is used to find the minimum? (Heap Sort)

Outline

- - Comparision
- **Bubble Sort**
 - Algo

 - Analysis
- Insertion Sort
 - Algo
 - Example
 - Analysis
- - Algo

 - Analysis
- Merge Sort
 - Algo



- Divide the *n*-elements sequence into 2 subsequences of n/2 elements.
- Recursively sort each subsequence using merge sort.
- Merge the two sorted subsequences to produce the sorted answer.





- Divide the *n*-elements sequence into 2 subsequences of n/2 elements.
- Recursively sort each subsequence using merge sort.
- Merge the two sorted subsequences to produce the sorted answer.



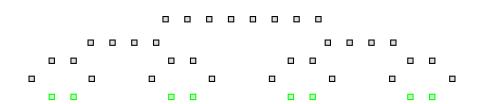
- Divide the *n*-elements sequence into 2 subsequences of n/2 elements.
- Recursively sort each subsequence using merge sort.
- Merge the two sorted subsequences to produce the sorted answer.



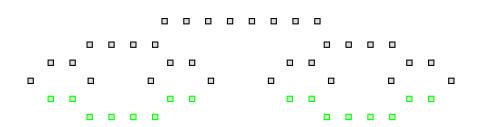
- Divide the *n*-elements sequence into 2 subsequences of n/2 elements.
- Recursively sort each subsequence using merge sort.
- Merge the two sorted subsequences to produce the sorted answer.



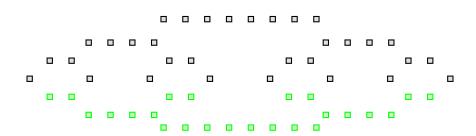
- Divide the *n*-elements sequence into 2 subsequences of n/2 elements.
- Recursively sort each subsequence using merge sort.
- Merge the two sorted subsequences to produce the sorted answer.

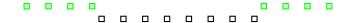


- Divide the *n*-elements sequence into 2 subsequences of n/2 elements.
- Recursively sort each subsequence using merge sort.
- Merge the two sorted subsequences to produce the sorted answer.

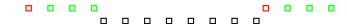


- Divide the *n*-elements sequence into 2 subsequences of n/2 elements.
- Recursively sort each subsequence using merge sort.
- Merge the two sorted subsequences to produce the sorted answer.



































































5 1 2 0 4 6 0 9

5 1 2 0 4 6 0 9

5

2

0

4

6

0

5 1 2 0 4 6 0 9

5

1

2

0

4

6

0

9

Pass 1

5 1 2 0 4 6 0 9

5

1

2

0

4

6

0

9

Pass 1

1 5

5 1 2 0 4 6 0 9

5

1 2

0

4

6

0

9

Pass 1

1 5

0 2

5 1 2 0 4 6 0 9

5 1 2 0 4 6 0 9

Pass 1

1 5 0 2 4 6 0 9

5 1 2 0 4 6 0 9

Pass 1

1 5

0 2

4

0 9

Pass 2

1 5

Pass 1

1 5

0 2

4

0 9

Pass 2

1 5

2

5 1 2 0 4 6 0 9

Pass 1

1 5

0 2

4 6

0 9

Pass 2

1 5

2

5 1 2 0 4 6 0 9

Pass 1

1 5

0 2

4

0 9

Pass 2

5

2

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5 1 2 0 4 6 0 9

Pass 1

1 5

0 2

4

0 9

Pass 2

5

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5 1 2 0 4 6 0 9

Pass 1

1 5

0 2

4

0 9

Pass 2

5

0 1 2 -

5 1 2 0 4 6 0 9

Pass 1

1 5

0 2

4

0 9

Pass 2

 0 1 2 5

5 1 2 0 4 6 0 9

Pass 1

1 5

0 2

4

0 9

Pass 2

4

0 9

5 1 2 0 4 6 0 9

Pass 1

1 5

0 2

4 6

0 9

Pass 2

5 0 1 2

0 4 6 9

5 1 2 0 4 6 0 9

Pass 1

1 5

0 2

4

0 9

Pass 2

0 1 2 5

0 4 6 9

Pass 3

0 1 2 5

0 4 6 9

5 1 2 0 4 6 0 9

Pass 1

1 5

0 2

4 6

0 9

Pass 2

5 0 1 2

4 6 9











5 1 2 0 4 6 0 9



Algorithm 16 MergeSort algorithm

```
1: procedure MERGESORT(a[], left, right)
2:
       if left < right then
3:
           mid \leftarrow |(left + right)/2|
           MergeSort(a, left, mid)
5:
           MergeSort(a, mid + 1, right)
6:
           Merge(a, left, mid, right)
       end if
8: end procedure
9: procedure MERGE(a[], left, mid, right)
10:
        L \leftarrow a[left..mid]
11:
        R \leftarrow a[mid + 1..right]
12:
        k \leftarrow left
13:
        while there are still elements in L or R do
14:
            Compare the 'first' elements in L and R
15:
           Move the minimum of them from its corresponding list to a[k]
16:
            k \leftarrow k + 1
17:
        end while
18: end procedure
```

Time analysis

- If the problem size is small, say $n \le c$ for some constant c, we can solve the problem in constant time, i.e., $\Theta(1)$.
- Let T(n) be the time needed to sort for input of size n.
- Let C(n) be the time needed to merge 2 lists of total size n. We know that $C(n) = \Theta(n)$.
- Assume that the problem can be split into 2 subproblems in constant time and c = 1, then

$$T(n) = \begin{cases} \Theta(1) & \text{if } n \leq 1 \\ 2T(\frac{n}{2}) + \Theta(n) & \text{if } n > 1. \end{cases}$$



$$T(n) = 2T(\frac{n}{2}) + cn$$

$$= 2[2T(\frac{n}{4}) + c\frac{n}{2}] + cn$$

$$= 4T(\frac{n}{4}) + 2cn$$

$$= 4[2T(\frac{n}{8}) + c\frac{n}{4}] + 2cn$$

$$= 8T(\frac{n}{8}) + 3cn$$

$$= \vdots$$

$$= 2^k T(\frac{n}{2^k}) + kcn$$

$$w.l.o.g., assume $2^k = n \Rightarrow k = \log_2 n$

$$T(n) = cn\log_2 n + n = \Theta(n\log_2 n)$$$$

- Worst case:
- Best case:
- Reverse sorted:
- Already sorted:

In-place

Stable

- Worst case: O(n log₂ n)
- Best case: O(n log₂ n)
- Reverse sorted:
- Already sorted:

In-place

Stable

- Worst case: O(n log₂ n)
- Best case: $O(n \log_2 n)$
- Reverse sorted: $\Theta(n \log_2 n)$
- Already sorted:

In-place

Stable

- Worst case: O(n log₂ n)
- Best case: O(n log₂ n)
- Reverse sorted: $\Theta(n \log_2 n)$
- Already sorted: Current implementation: $\Theta(n \log_2 n)$. Can be made linear. Any suggestion?

In-place

Stable



- Worst case: O(n log₂ n)
- Best case: $O(n \log_2 n)$
- Reverse sorted: $\Theta(n \log_2 n)$
- Already sorted: Current implementation: $\Theta(n \log_2 n)$. Can be made linear. Any suggestion?

In-place

Stable



- Worst case: O(n log₂ n)
- Best case: O(n log₂ n)
- Reverse sorted: $\Theta(n \log_2 n)$
- Already sorted: Current implementation: $\Theta(n \log_2 n)$. Can be made linear. Any suggestion?

In-place

O(1) extra space

Stable



- Worst case: O(n log₂ n)
- Best case: O(n log₂ n)
- Reverse sorted: $\Theta(n \log_2 n)$
- Already sorted: Current implementation: $\Theta(n \log_2 n)$. Can be made linear. Any suggestion?

In-place

O(1) extra space

Stable

Yes



- Worst case: O(n log₂ n)
- Best case: O(n log₂ n)
- Reverse sorted: $\Theta(n \log_2 n)$
- Already sorted: Current implementation: $\Theta(n \log_2 n)$. Can be made linear. Any suggestion?

In-place

O(1) extra space

Stable

Yes

Adaptive

No (Current implementation)/ Yes (After a modification)

References

- Cormen, Leiserson, Rivest: Introduction to Algorithms
- http://www.sorting-algorithms.com/

